

Dalarna

**A Simplistic Capability-Based Dynamic
Language Design For Data Race Freedom**



Kiko Fernandez Reyes¹

James Noble²

Isaac Oscar Gariano²

Erin Greenwood-Thressman²

Michael Homer²

Tobias Wrigstad¹

¹Uppsala University

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Motivation

- Object-oriented dynamic languages in the rise [1]:



- All programming languages have concurrency constructs

Goroutines

Async & Await

Task

Web Workers

Promise & Future

[1] Most Loved Language from Stackoverflow'20 Survey:

<https://insights.stackoverflow.com/survey/2020#technology-most-loved-dreaded-and-wanted-languages-loved>

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Legend

	Dynamic
	Unsafe (not data race free)
	Deep copy

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Dalarna Programming Model



- Focused on *dynamic* languages
- Make programs *data race free* by mapping objects to *capabilities*:
 - Remove *data races* in   ( )
 - Remove *deep copying* where possible  
- Embedded Dalarna into the Grace language

Dalarna Programming Model

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<i>Objects</i>	O	$::=$	$K \mathbf{obj} \{\overline{F} \overline{M}\}$
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Dalarna Programming Model

- Capabilities annotations

Object	Effects			
	Read	Write	Alias	Transfer
Imm	•	×	•	•
Iso	•	•	×	•
Local	•	•	•	×
Unsafe	•	•	•	•

Dalarna Programming Model

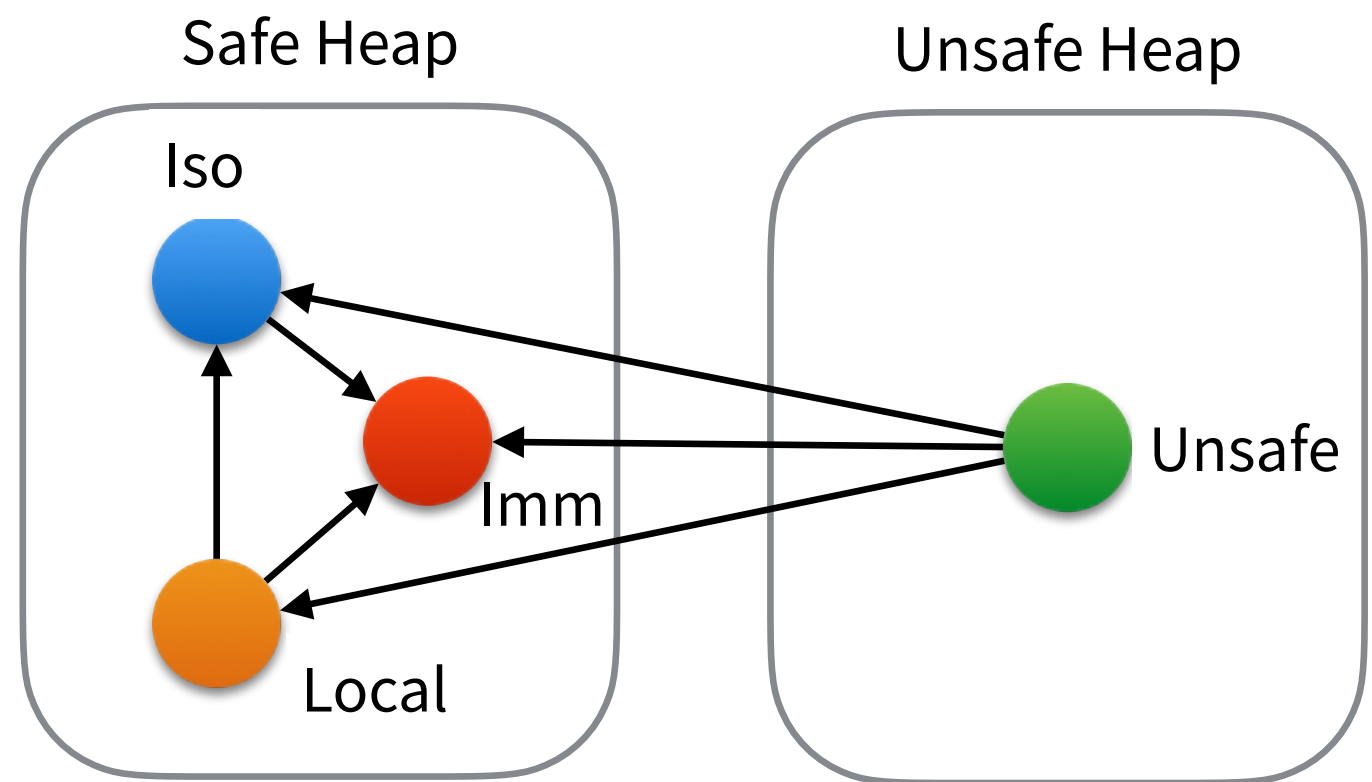
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Object	Effects			
	Read	Write	Alias	Transfer
Imm	•	×	•	•
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Local & Imm = Read and Alias

Dalarna Programming Model

- Object capability containment relation



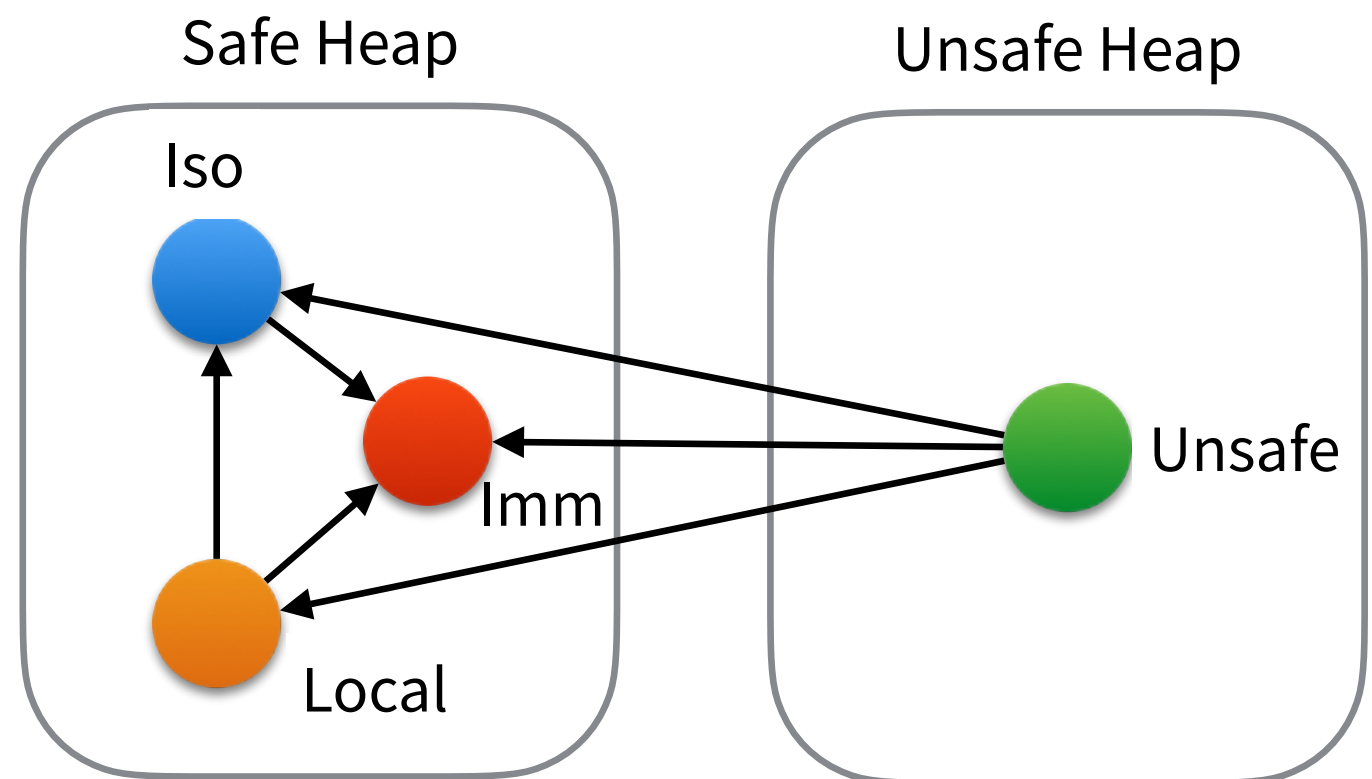
Dalarna Programming Model

- Object capability containment relation

This relation must **hold**:

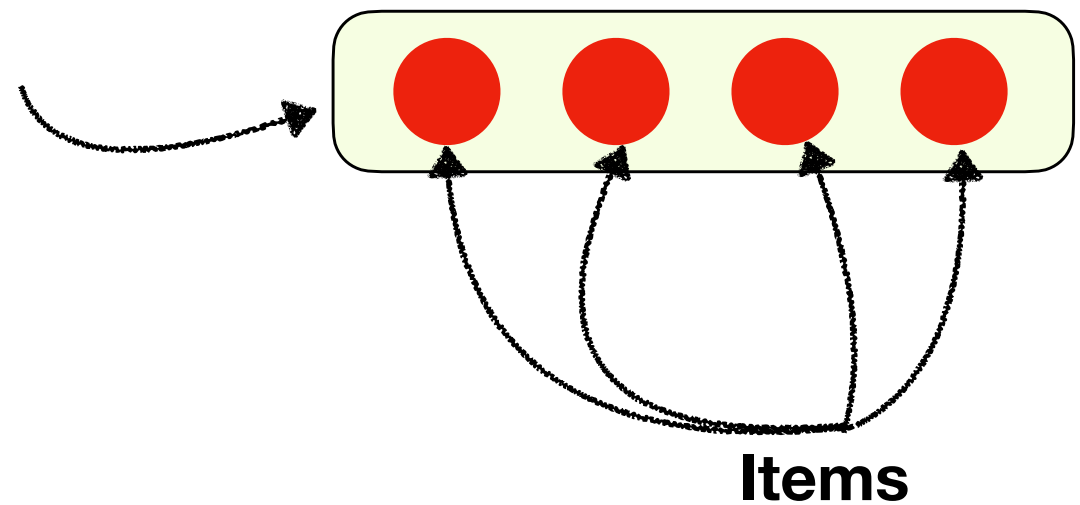
$\text{unsafe} \leq \text{local} \leq \text{iso} \leq \text{imm}$

(Reflexive & Transitive)



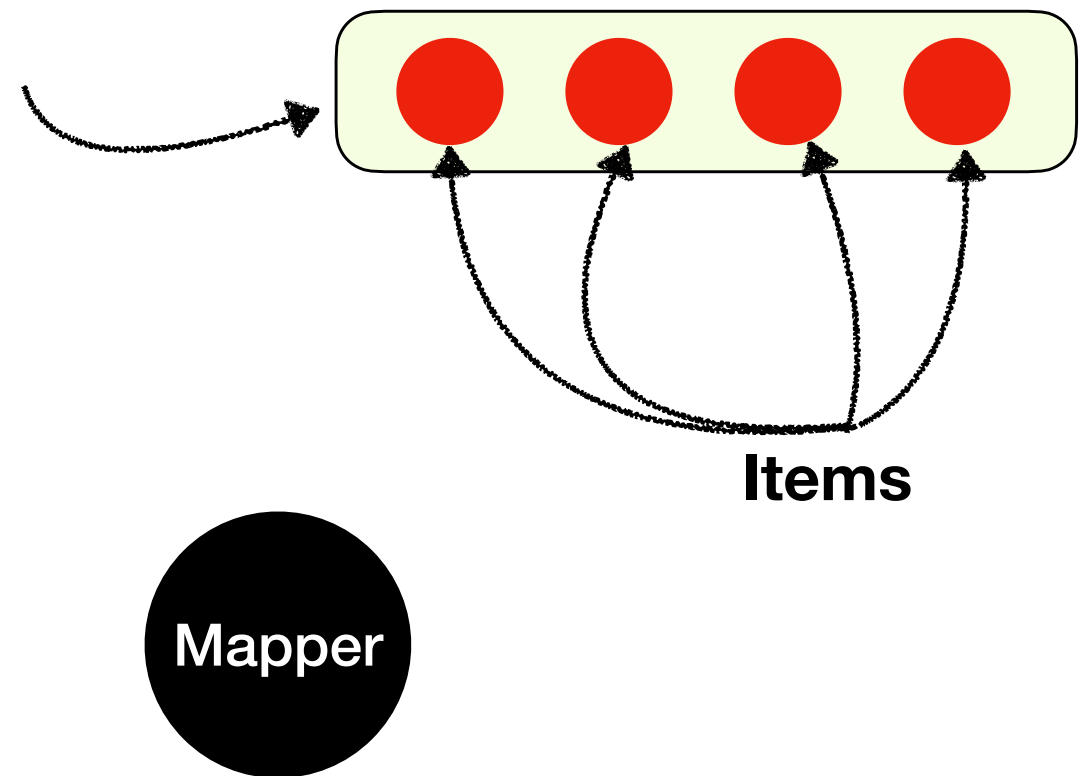
Dalarna Programming Model

Initial code



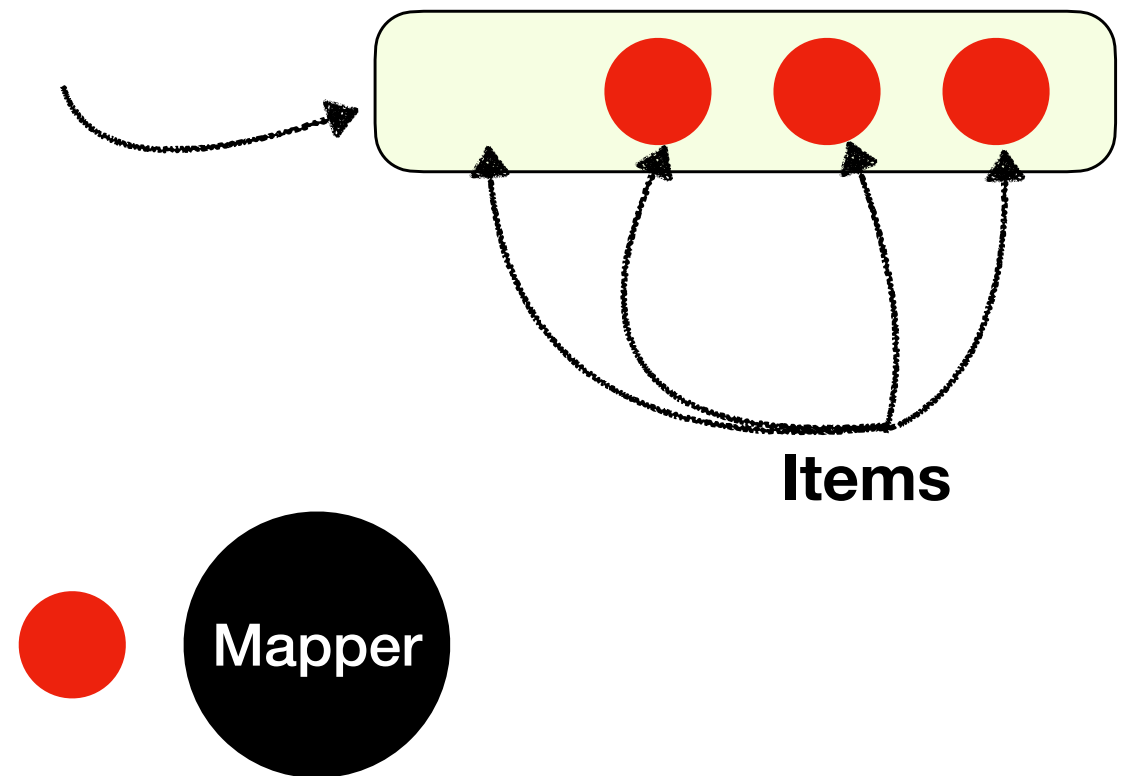
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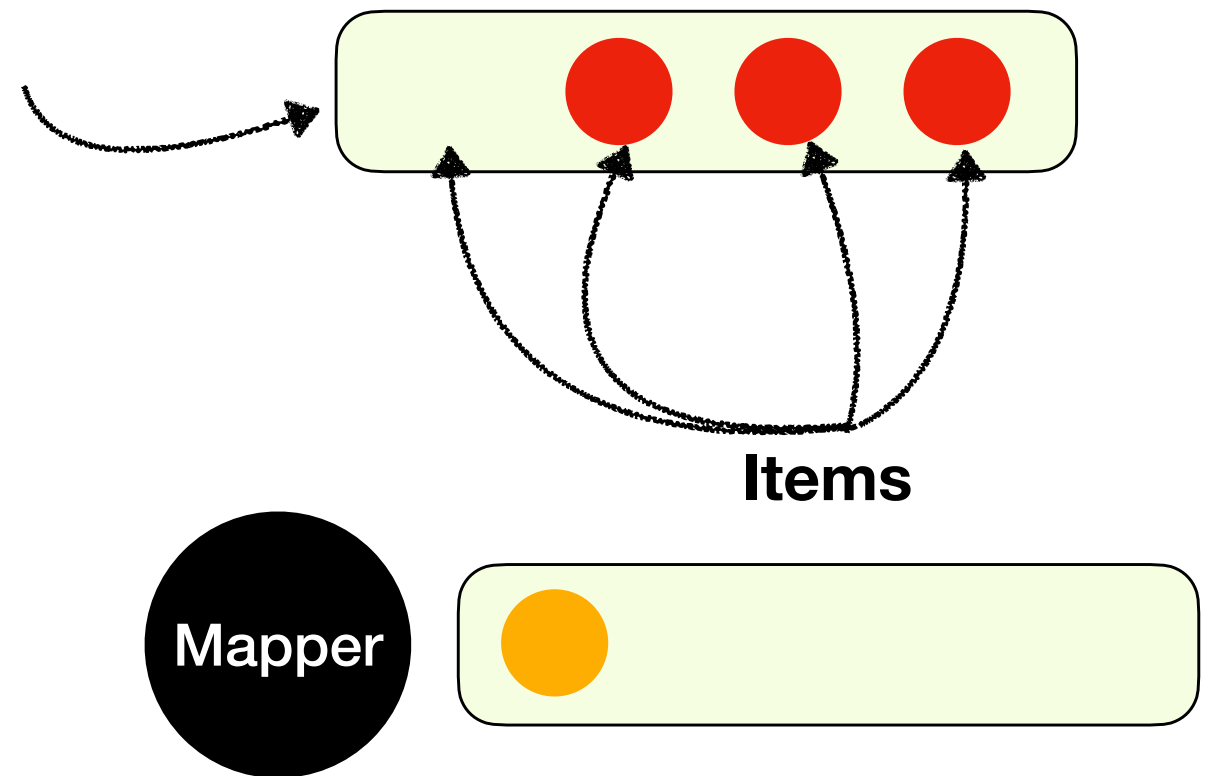
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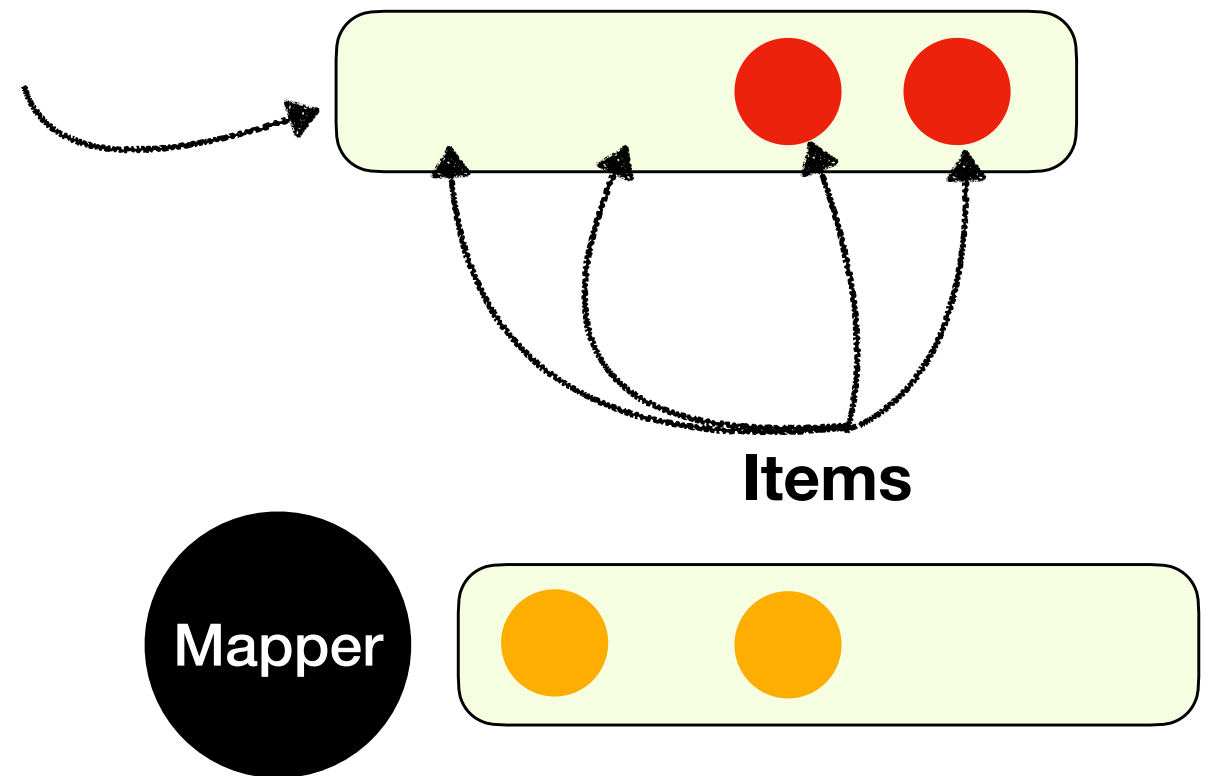
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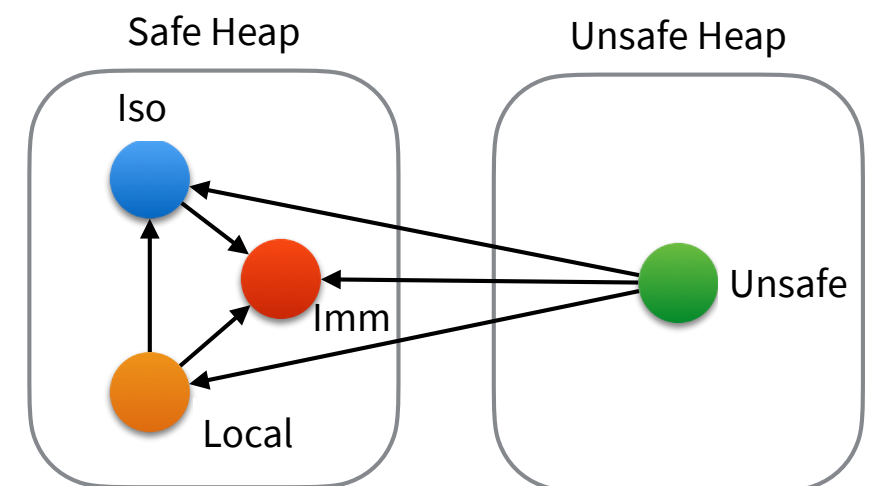
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ls.append(object 0 {...} )

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    let ch = spawn(ch) {
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    }
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ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}
≡
ls = unsafe object List {...}
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Dalarna Programming Model

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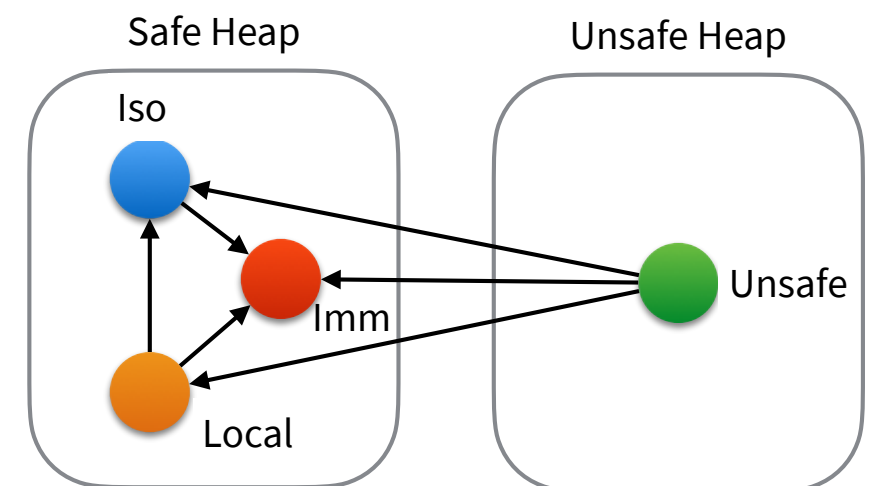
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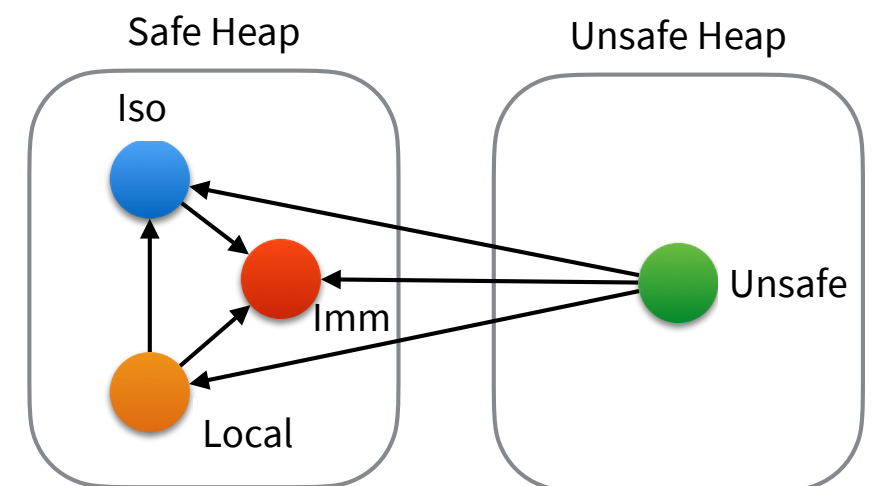
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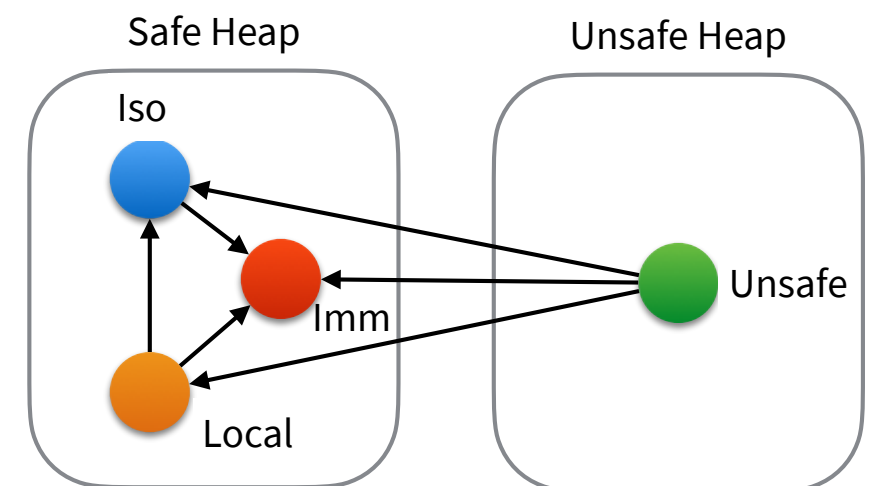
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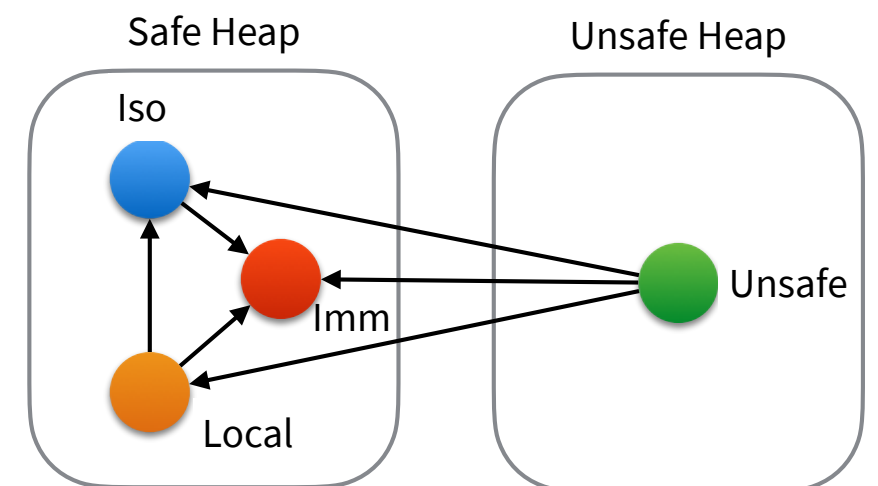
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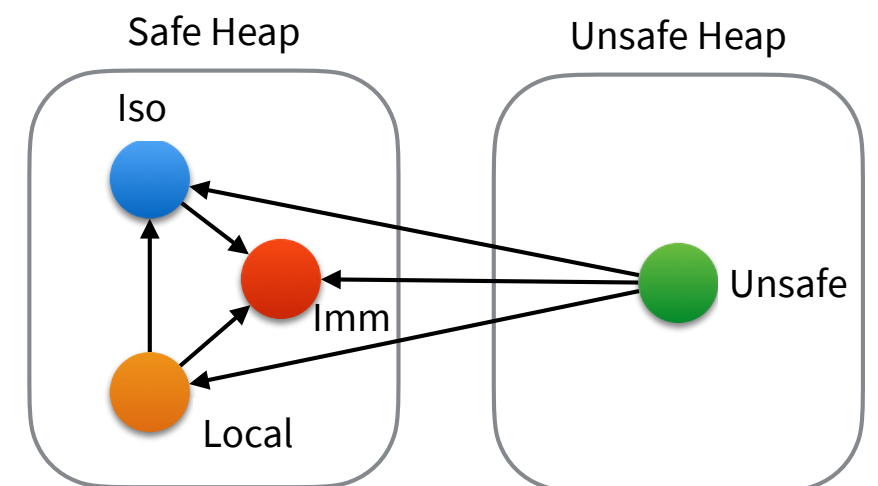
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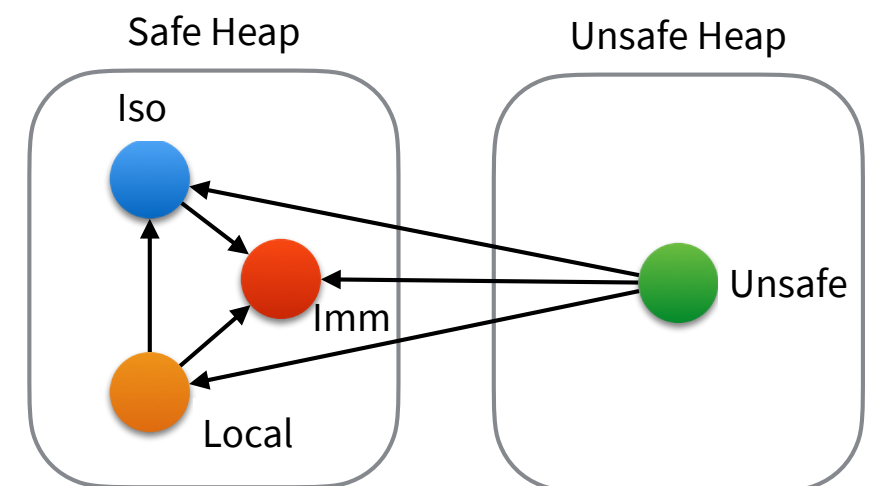
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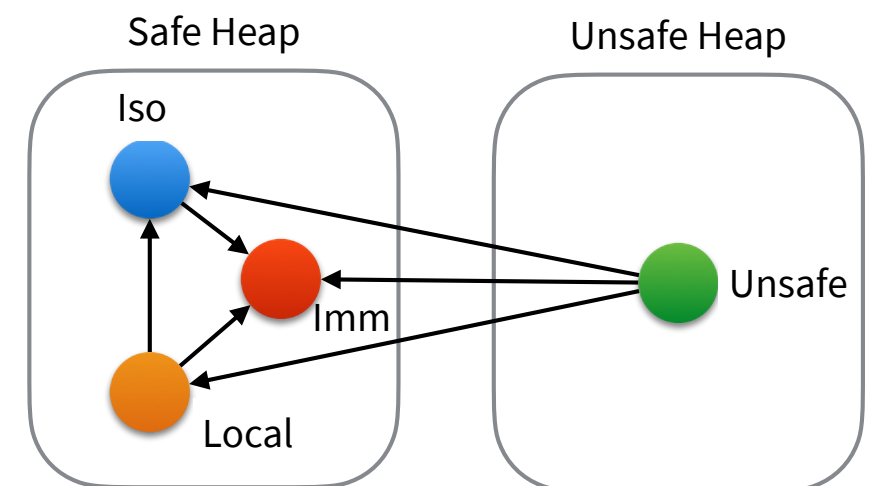
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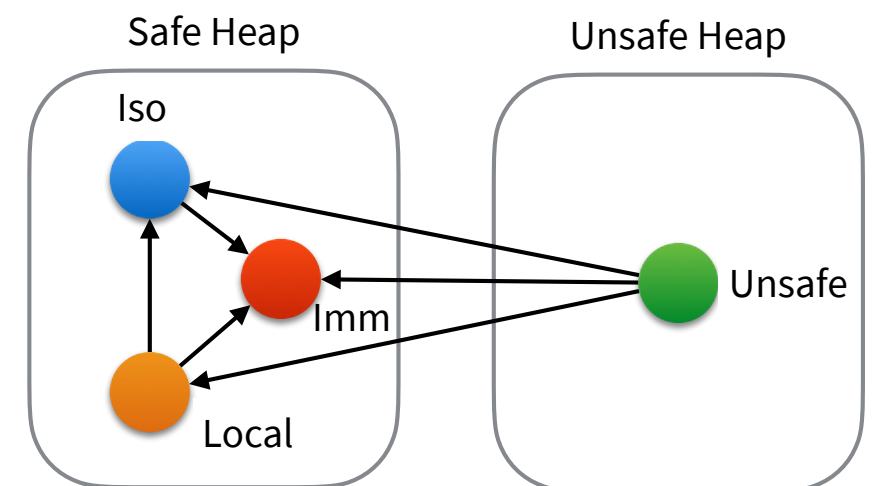
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ls.map(mapper) This code is **not** data race free

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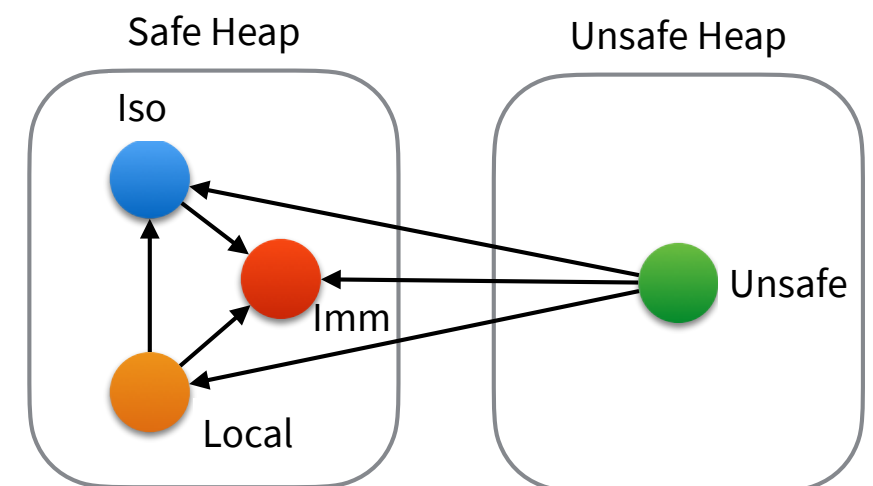
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This code is **not** data race free
Do not share state,
but there is no enforcement

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Dalarna Programming Model

~~Initial code~~ **Isolated list**

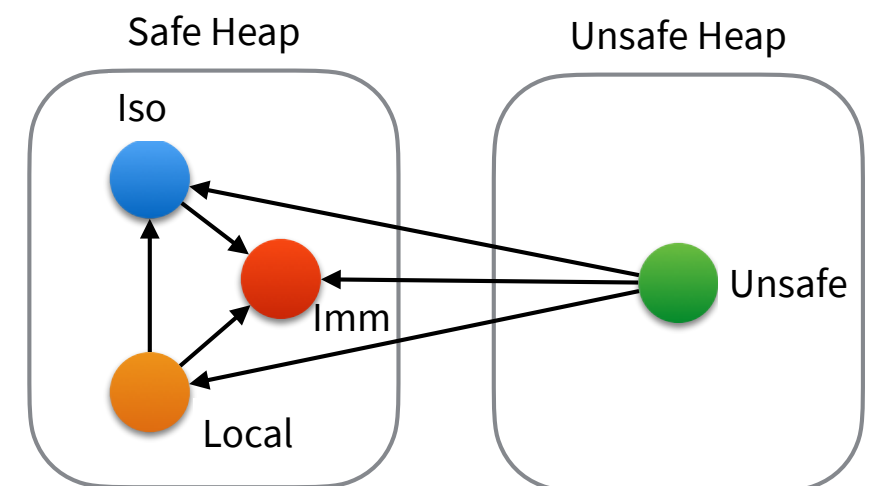
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    ch <- o
    o.f = ...  -- mutate object
  }
}

ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}
≡
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ **Isolate** list

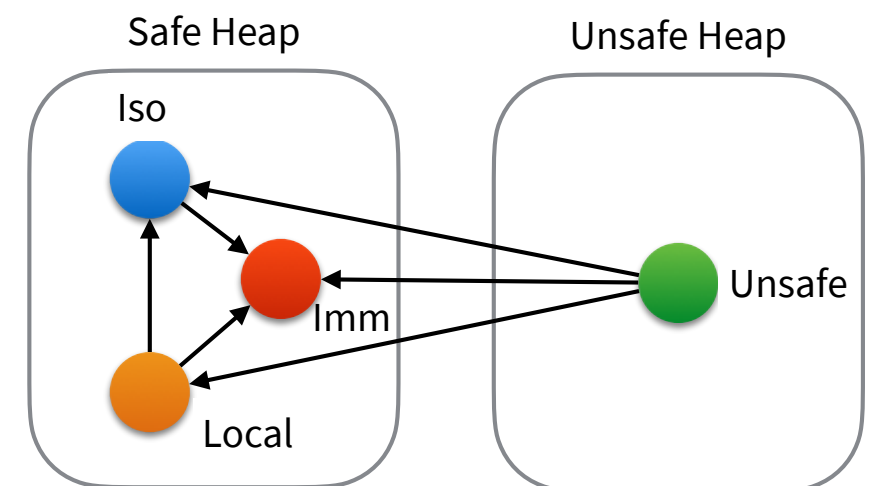
```
ls = iso object List {...}  
ls.append(object 0 {...} )
```

```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ...    -- mutate object  
    }  
    ch <- o  
    o.f = ...  -- mutate object  
  }  
}
```

```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ **Isolate** list

```
ls = iso object List {...}  
ls.append(object 0 {...})
```



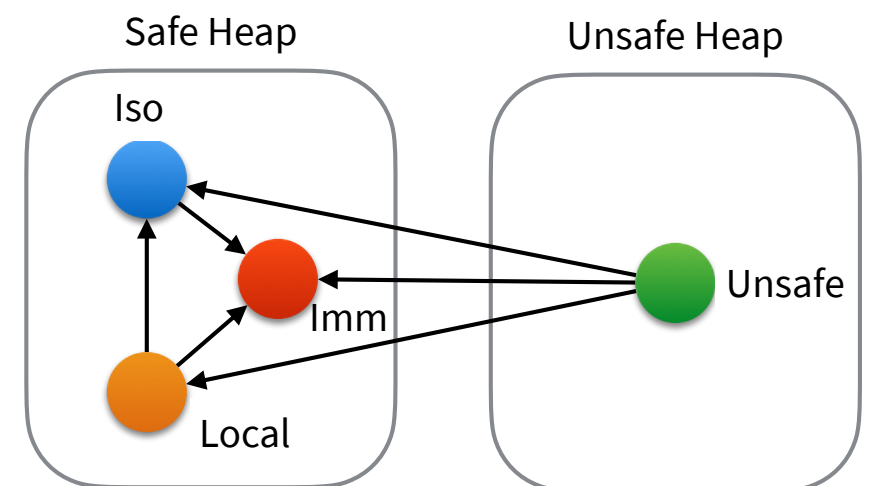
Broken containment relation

```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- o  
    o.f = ... -- mutate object  
  }  
}
```

```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ **Isolate** list

```
ls = iso object List {...}  
ls.append(imm object 0 {...})
```

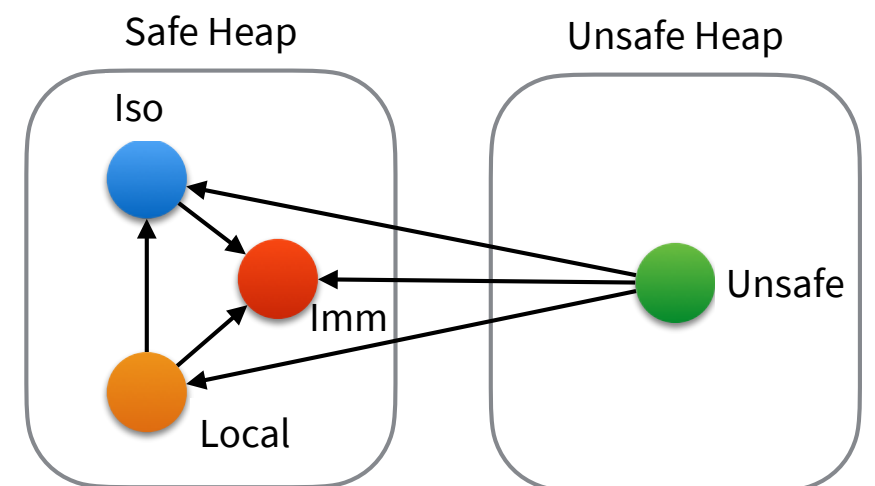


```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- o  
    o.f = ... -- mutate object  
  }  
}
```

```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ **Isolate** list

```
ls = iso object List {...}  
ls.append(imm object 0 {...})
```

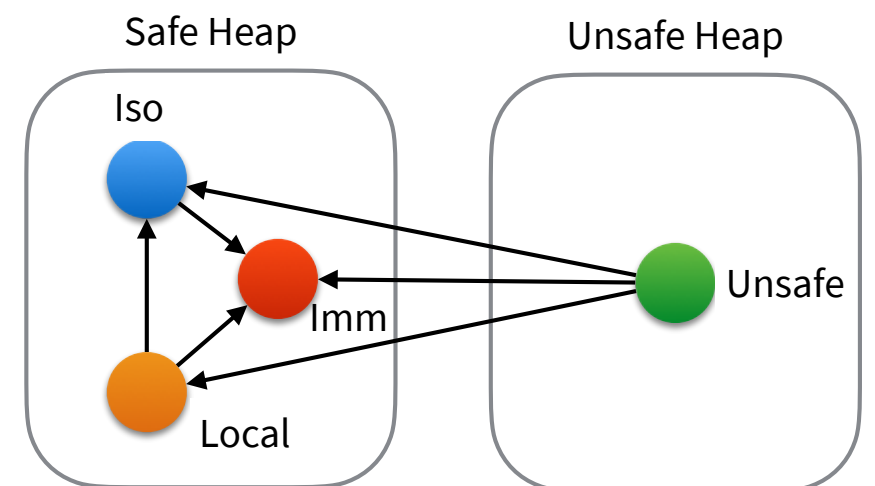


```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- o  
    o.f = ... -- mutate object  
  }  
}
```

```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ **Isolate** list

```
ls = iso object List {...}  
ls.append(imm object 0 {...})
```

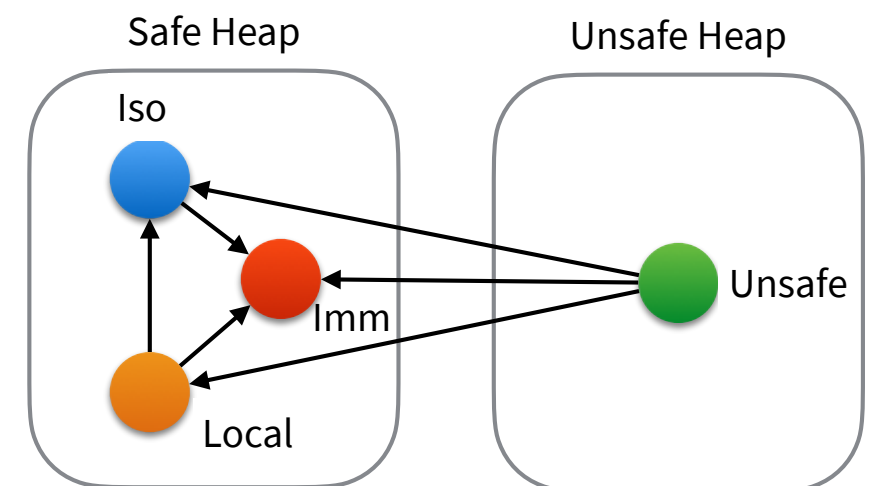


```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- o  
    o.f = ... -- mutate object  
  }  
}
```

```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ **Isolate** list

```
ls = iso object List {...}  
ls.append(imm object 0 {...})
```



```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- o  
    o.f = ... -- mutate object  
  }  
}
```

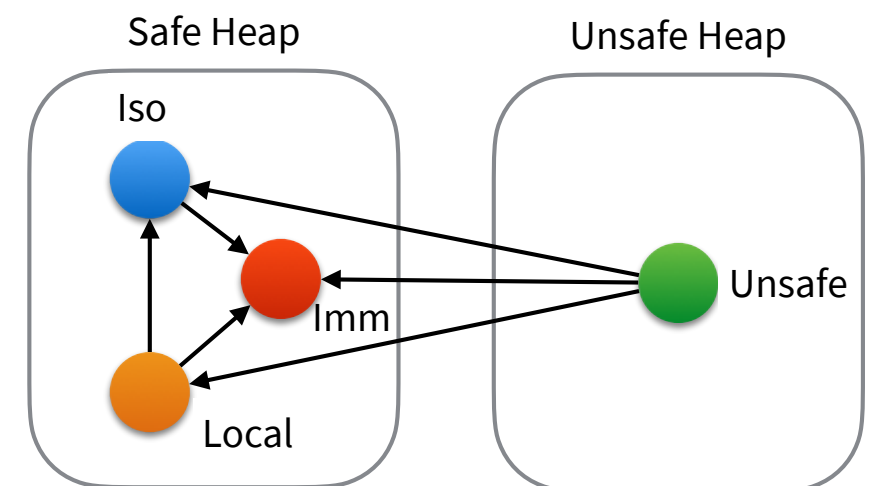
ls.map(mapper)

Zero copying



Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ **Isolate** list

```
ls = iso object List {...}  
ls.append(imm object 0 {...})
```

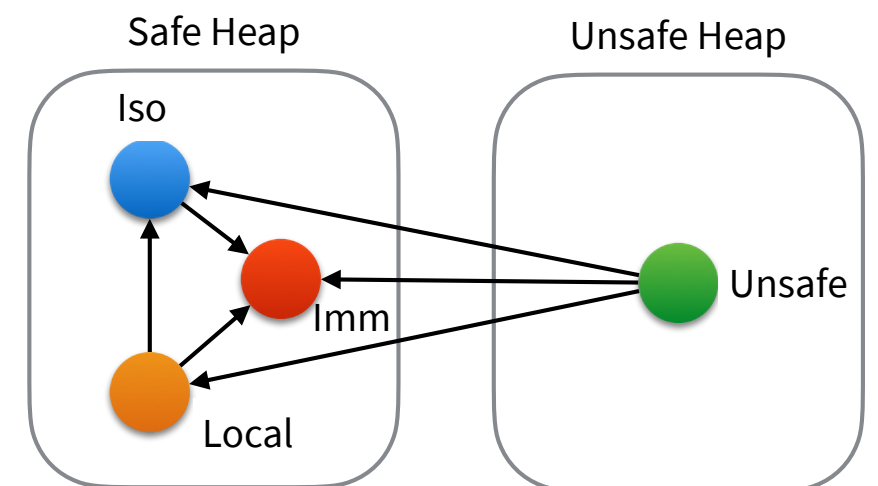


```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object X  
    }  
    ch <- o  
    o.f = ... -- mutate object X  
  }  
}
```

```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ **Isolate** list

```
ls = iso object List {...}  
ls.append(imm object 0 {...})
```



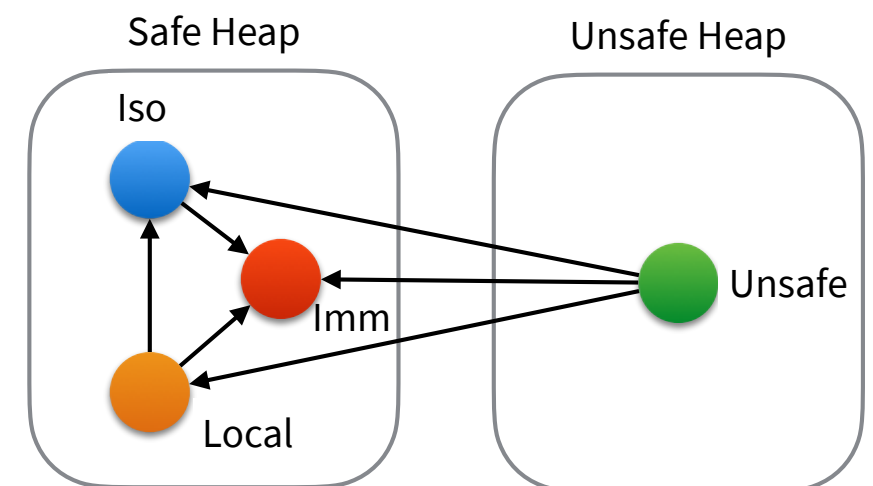
Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```

```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object X  
    }  
    ch <- o  
    o.f = ... -- mutate object X  
  }  
}
```

Immutable objects cannot be mutated

```
ls.map(mapper)
```



Dalarna Programming Model

~~Initial code~~ Local list

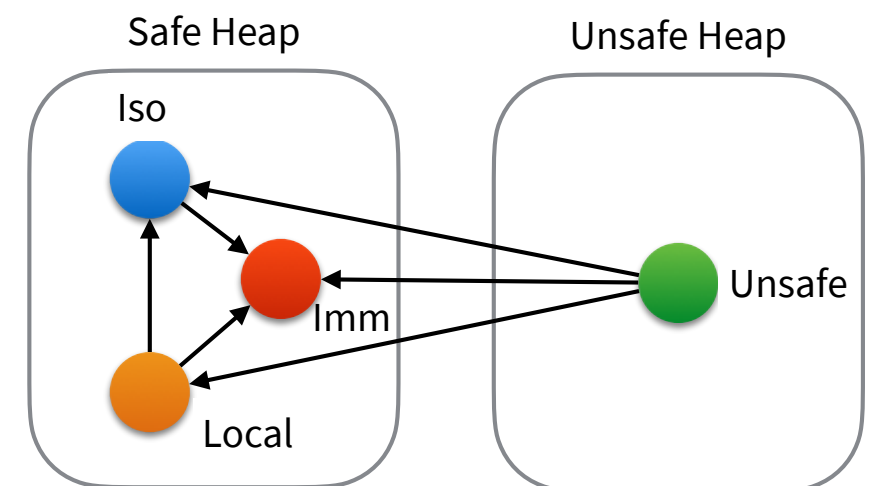
```
ls = local object List {...}  
ls.append(local object 0 {...} )
```

```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ...      -- mutate object  
    }  
    ch <- o  
    o.f = ...  -- mutate object  
  }  
}
```

```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = local object List {...}  
ls.append(local object 0 {...} )
```

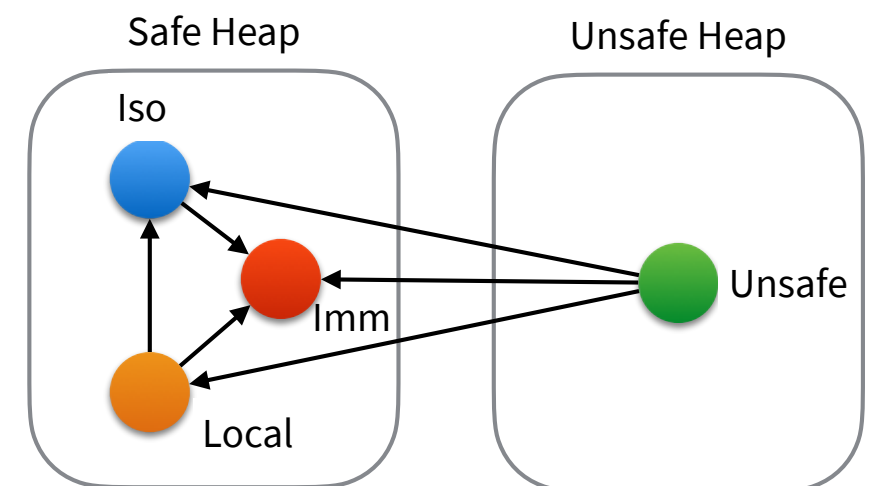


```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- o  
    o.f = ... -- mutate object  
  }  
}
```

```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = local object List {...}  
ls.append(local object 0 {...} )
```

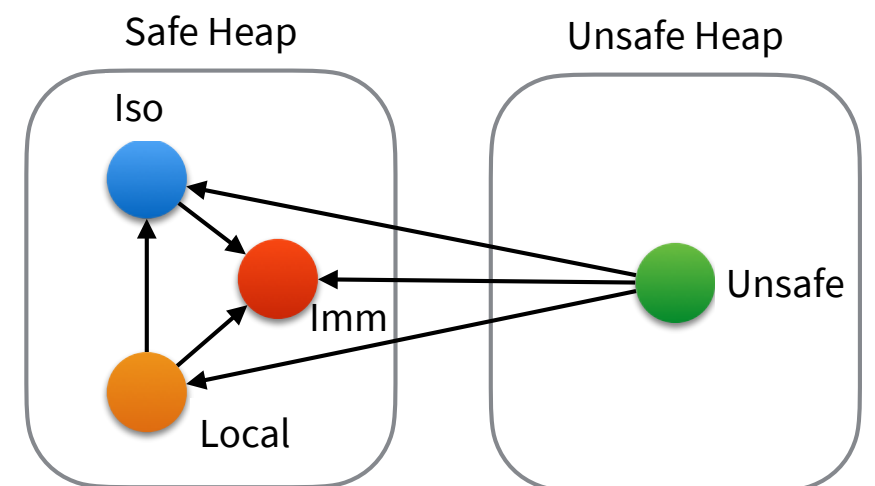


```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- o  
    o.f = ... -- mutate object  
  }  
}
```

```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = local object List {...}  
ls.append(local object 0 {...} )
```

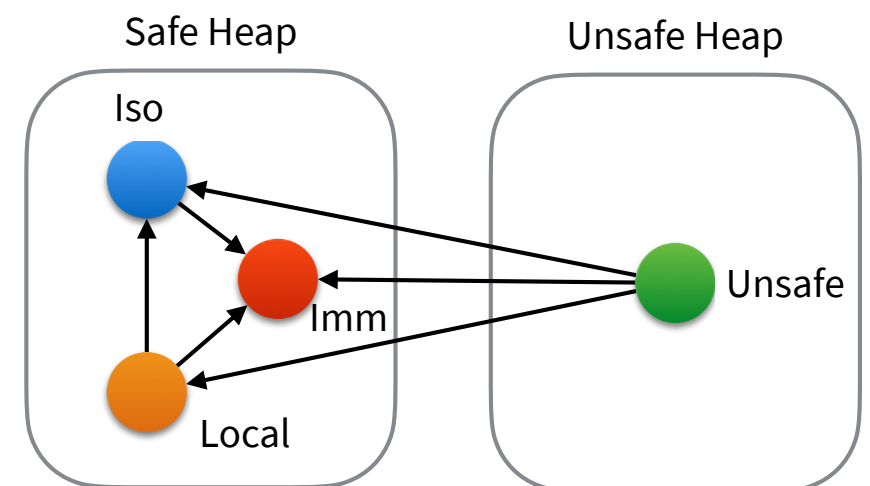


```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- o  
    o.f = ... -- mutate object  
  }  
}
```

```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = local object List {...}  
ls.append(local object 0 {...} )
```



Default object considered with
unsafe capability

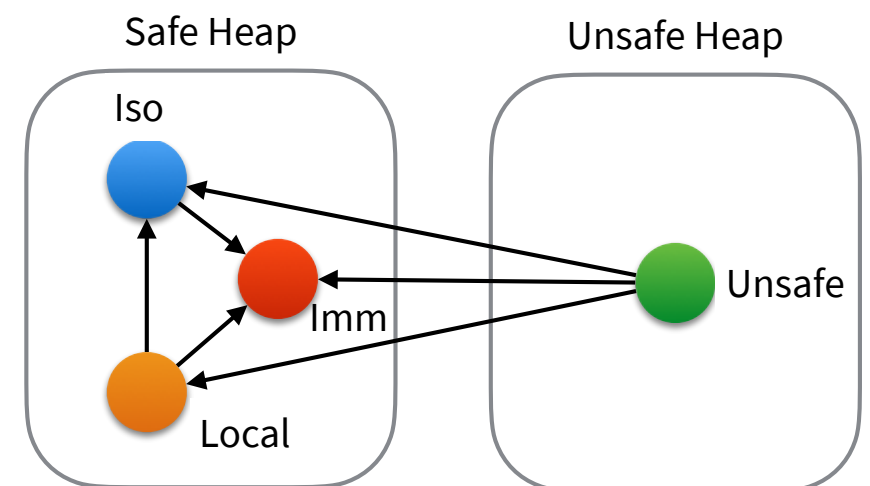
```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```

```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- o  
    o.f = ... -- mutate object  
  }  
}
```



Local objects cannot move across threads

```
ls.map(mapper)
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = local object List {...}  
ls.append(local object 0 {...} )
```



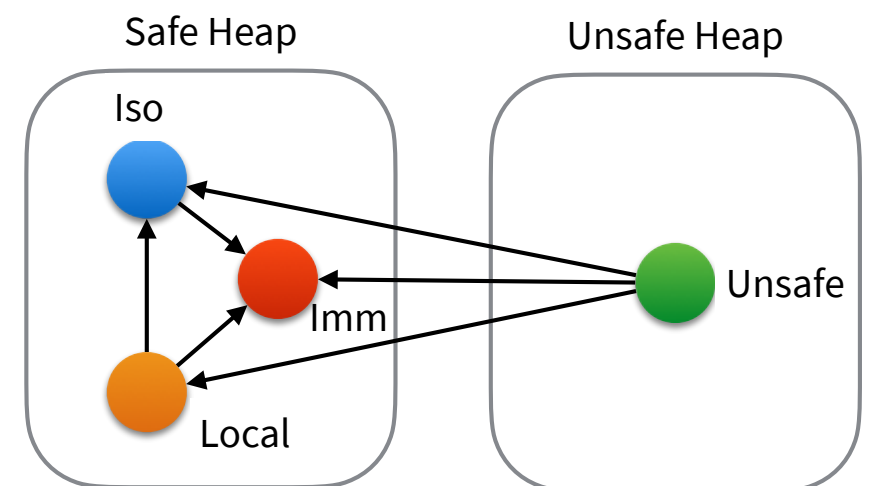
```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- copy(o)  
    o.f = ... -- mutate object  
  }  
}
```



```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = local object List {...}  
ls.append(local object 0 {...} )
```



```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- copy(o)  
    o.f = ... -- mutate object  
  }  
}
```



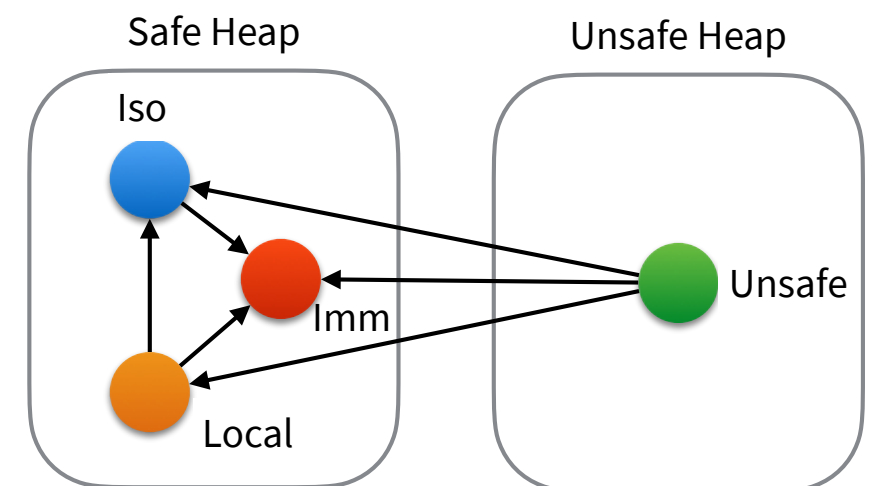
TypeScript



```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = local object List {...}  
ls.append(local object 0 {...} )
```



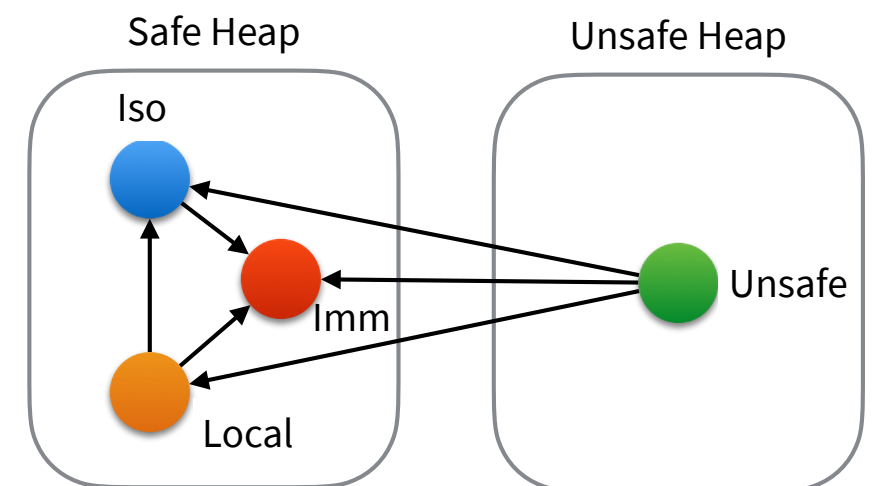
```
mapper = object Mapper(...) {  
  method map(o) {  
    let ch = spawn(ch) {  
      oo <- ch  
      ... -- mutate object  
    }  
    ch <- unsafe copy(o)  
    o.f = ... -- mutate object  
  }  
}
```



```
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Unsafe Local list

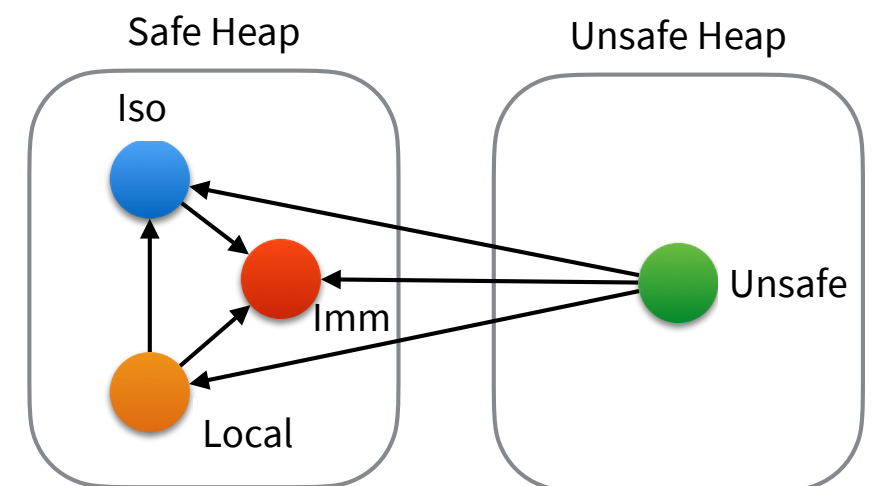
```
ls = unsafe object List {...}  
ls.append(local object 0 {...} )
```

```
let ch = spawn(ch) {  
  oo <- ch  
  ...    -- mutate object  
}
```

```
ch <- ls
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Unsafe Local list

```
ls = unsafe object List {...}  
ls.append(local object 0 {...} )
```

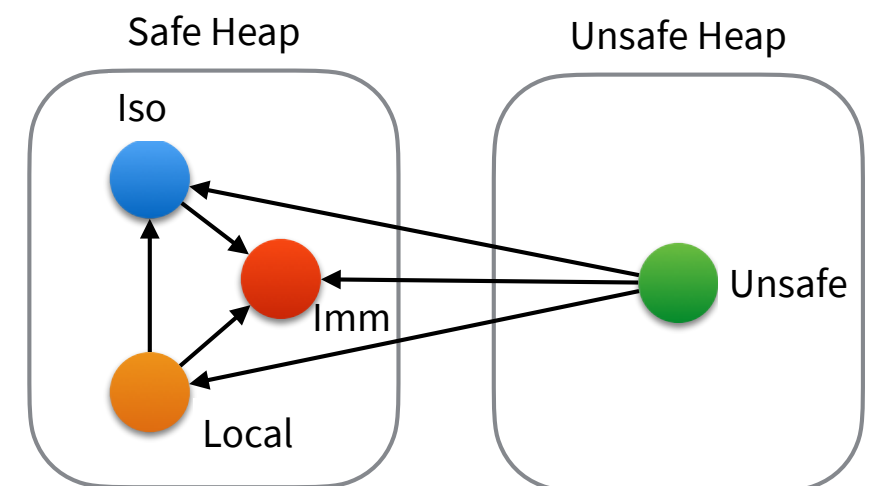


```
let ch = spawn(ch) {  
  oo <- ch  
  ...    -- mutate object  
}
```

```
ch <- ls
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = unsafe object List {...}  
ls.append(local object 0 {...} )
```

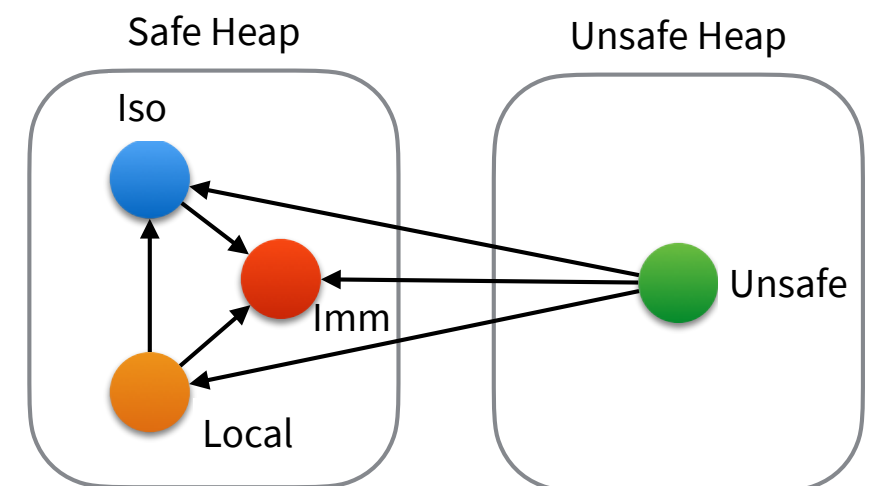


```
let ch = spawn(ch) {  
  oo <- ch  
  ...    -- mutate object  
}
```

```
ch <- ls
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = unsafe object List {...}  
ls.append(local object 0 {...} )
```

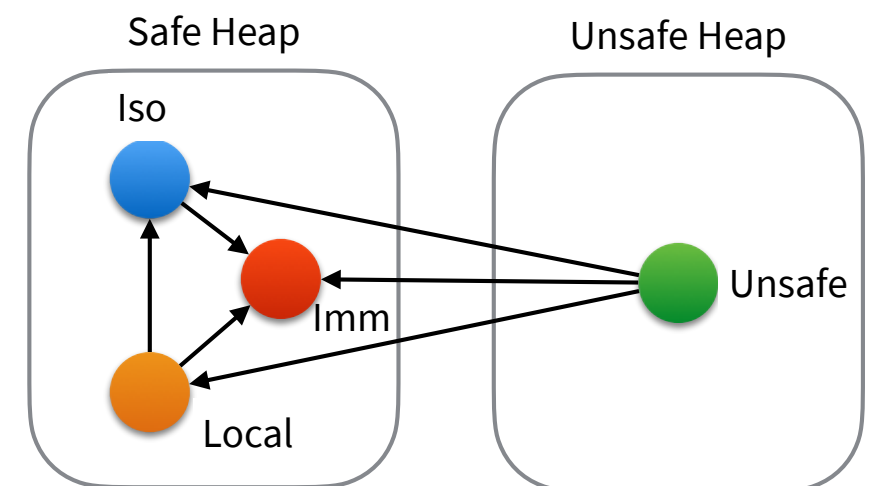


```
let ch = spawn(ch) {  
  oo <- ch  
  ...    -- mutate object  
}
```

```
ch <- ls
```

Default object considered with
unsafe capability

```
ls = object List {...}  
≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = unsafe object List {...}  
ls.append(local object 0 {...} )
```



```
let ch = spawn(ch) {  
  oo <- ch  
  ...    -- mutate object  
}
```

```
ch <- ls
```



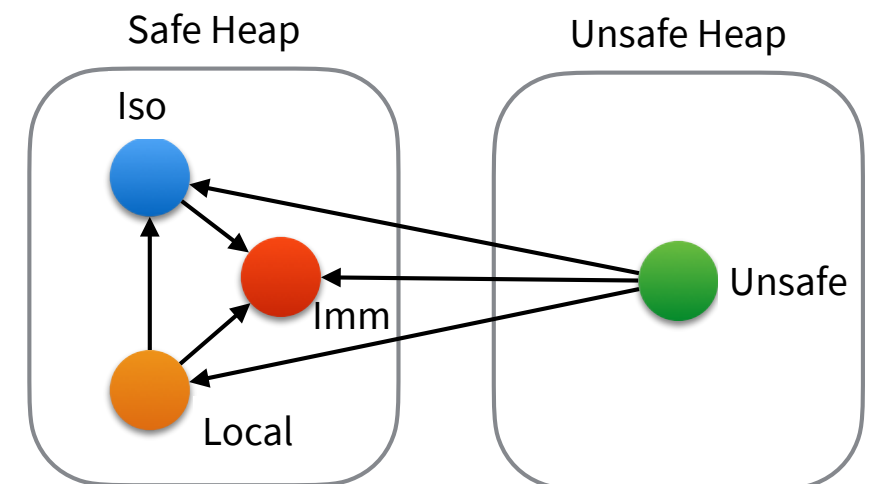
Local objects cannot move across threads

Default object considered with
unsafe capability

```
ls = object List {...}
```

≡

```
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = unsafe object List {...}  
ls.append(local object 0 {...} )
```



```
let ch = spawn(ch) {  
  oo <- ch  
  ...    -- mutate object  
}
```

```
ch <- ls
```



Local objects cannot move across threads

Main Point:

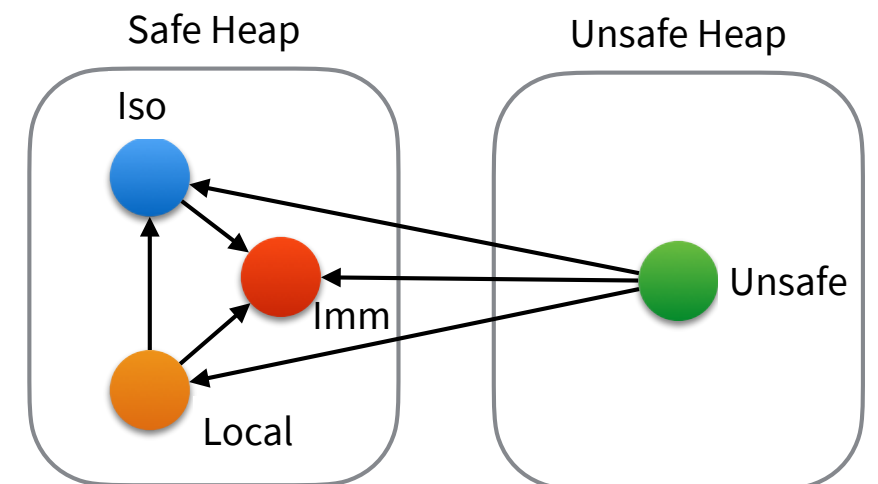
*Any code that can lead
to a data race is a runtime error*

Default object considered with
unsafe capability

```
ls = object List {...}
```

≡

```
ls = unsafe object List {...}
```



Dalarna Dynamic Semantics

Configuration Cfg $::=$ $H, \bar{T}$
 Thread T $::=$ $t \mid Err$
 Thread Err Err $::=$ $Err_N \mid Err_A \mid Err_P \mid Err_C$

Dalarna Dynamic Semantics

$Configuration \quad Cfg ::= H, \bar{T}$
 $Thread \quad T ::= t \mid Err$
 $Thread \quad Err ::= Err_N \mid Err_A \mid Err_P \mid Err_C$

$\frac{(R-LET) \quad x \notin dom(H)}{H; \mathbf{let} \ x = v \ \mathbf{in} \ t \rightsquigarrow H, x \mapsto v; t}$	$\frac{(R-VAR) \quad H(x) = v \quad \neg isIso(H, v)}{H; E[x] \rightsquigarrow H; E[v]}$	$\frac{(R-CONSUME) \quad H(x) = v \quad H' = H[x \mapsto \top]}{H; E[\mathbf{consume} \ x] \rightsquigarrow H'; E[v]}$	$\frac{(R-FIELD) \quad H(H(x)) = _ \mathbf{obj} \ \{ _ f = v \ \overline{M} \} \quad \neg isIso(H, v)}{H; E[x.f] \rightsquigarrow H; E[v]}$
$\frac{(R-UPDATE) \quad H(x) = \iota \quad H(\iota) = K \ \mathbf{obj} \ \{ \overline{f = v} \ f = v' \ \overline{M} \} \quad \neg isImm(H, \iota) \quad OkRef(H, K, v) \quad H' = H[\iota \mapsto K \ \mathbf{obj} \ \{ \overline{f = v} \ f = v' \ \overline{M} \}]}{H; E[x.f = v] \rightsquigarrow H'; E[v']}$	$\frac{(R-CASTLOC) \quad H(\iota) = K \ \mathbf{obj} \ \{ _ _ \}}{H; E[(K) \ \iota] \rightsquigarrow H; E[\iota]}$	$\frac{(R-NEW) \quad \forall f = v \in \overline{f = v}. OkRef(H, K, v) \quad \iota \text{ fresh} \quad H' = H, \iota \mapsto K \ \mathbf{obj} \ \{ \overline{f = v} \ \overline{M} \}}{H; E[K \ \mathbf{obj} \ \{ \overline{f = v} \ \overline{M} \}] \rightsquigarrow H'; E[\iota]}$	$\frac{(R-COPY) \quad \mathbf{iso} \notin K \quad OkDup(H, K, H(x)) = (H', \iota)}{H; E[K \ \mathbf{copy} \ x] \rightsquigarrow H'; E[\iota]}$
$\frac{(R-CALL) \quad x', y' \text{ fresh} \quad H(x) = \iota \quad H(\iota) = _ \mathbf{obj} \ \{ _ \overline{M} \ \mathbf{method} \ m(y) \ \{ t \} \}}{H; E[x.m(v)] \rightsquigarrow H, x' \mapsto \iota, y' \mapsto v; E[t[\mathbf{self} = x']][y = y']}]$	$\frac{(R-SPAWN) \quad \iota, i \text{ fresh} \quad x \notin dom(H) \quad H' = H, x \mapsto \iota, \iota \mapsto \mathbf{chan} \ \{ i, \emptyset \}}{H; E[\mathbf{spawn} \ (x) \ \{ t \}] \rightsquigarrow H'; E[\iota] \ t}$	$\frac{(R-RECV) \quad H(\iota) = \mathbf{chan} \ \{ i, \iota' \} \quad H' = H[\iota \mapsto \mathbf{chan} \ \{ i, \emptyset \}]}{H; E[\leftarrow \iota] \rightsquigarrow H'; E[\iota']}$	
$\frac{(R-SENDERBLOCK) \quad H(\iota) = \mathbf{chan} \ \{ _ , \emptyset \} \quad OkSend(H, v) \quad i \text{ fresh} \quad H' = H[\iota \mapsto \mathbf{chan} \ \{ i, v \}]}{H; E[\iota \leftarrow v] \rightsquigarrow H'; E[\blacksquare_i \ \iota]}$	$\frac{(R-SENDERUNBLOCK) \quad H(\iota) = \mathbf{chan} \ \{ i', v \} \quad v = \emptyset \vee i \neq i'}{H; E[\blacksquare_i \ \iota] \rightsquigarrow H; E[\iota]}$	$\frac{(REFCHECK) \quad H(\iota) = K' \ \mathbf{obj} \ \{ _ _ \} \quad OkField(K, K')}{OkRef(H, K, \iota)}$	$\frac{(HELPER-OKSEND) \quad H(\iota) = K \ \mathbf{obj} \ \{ _ _ \} \quad \neg \exists \iota' \in ROG(H, \iota) . isLocal(H, \iota')}{OkSend(H, \iota)}$
			$\frac{(HELPER-OKFIELD) \quad K_3 = Max(K_1) \quad \forall K_4 \in K_2 . K_3 \leq K_4}{OkField(K_1, K_2)}$

Dalarna Dynamic Semantics

$Configuration \quad Cfg ::= H, \bar{T}$
 $Thread \quad T ::= t \mid Err$
 $Thread \quad Err ::= Err_N \mid Err_A \mid Err_P \mid Err_C$

$\frac{\text{(R-LET)} \quad x \notin \text{dom}(H)}{H; \mathbf{let} \ x = v \ \mathbf{in} \ t \rightsquigarrow H, x \mapsto v; t}$	$\frac{\text{(R-VAR)} \quad H(x) = v \quad \neg \text{isIso}(H, v)}{H; E[x] \rightsquigarrow H; E[v]}$	$\frac{\text{(R-CONSUME)} \quad H(x) = v \quad H' = H[x \mapsto \top]}{H; E[\mathbf{consume} \ x] \rightsquigarrow H'; E[v]}$	$\frac{\text{(R-FIELD)} \quad H(H(x)) = _ \mathbf{obj} \ \{ _ f = v \ \overline{M} \} \quad \neg \text{isIso}(H, v)}{H; E[x.f] \rightsquigarrow H; E[v]}$	
$\frac{\text{(R-UPDATE)} \quad \begin{array}{l} H(x) = \iota \quad H(\iota) = K \ \mathbf{obj} \ \{ \overline{f = v} \ f = v' \ \overline{M} \} \\ \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\ H' = H[\iota \mapsto K \ \mathbf{obj} \ \{ \overline{f = v} \ f = v \ \overline{M} \}] \end{array}}{H; E[x.f = v] \rightsquigarrow H'; E[v']}$	$\frac{\text{(R-CASTLOC)} \quad H(\iota) = K \ \mathbf{obj} \ \{ _ _ \}}{H; E[(K) \ \iota] \rightsquigarrow H; E[\iota]}$	$\frac{\text{(R-NEW)} \quad \forall f = v \in \overline{f = v}. \text{OkRef}(H, K, v) \quad \iota \text{ fresh} \quad H' = H, \iota \mapsto K \ \mathbf{obj} \ \{ \overline{f = v} \ \overline{M} \}}{H; E[K \ \mathbf{obj} \ \{ \overline{f = v} \ \overline{M} \}] \rightsquigarrow H'; E[\iota]}$	$\frac{\text{(R-COPY)} \quad \mathbf{iso} \notin K \quad \text{OkDup}(H, K, H(x)) = (H', \iota)}{H; E[K \ \mathbf{copy} \ x] \rightsquigarrow H'; E[\iota]}$	
$\frac{\text{(R-CALL)} \quad \begin{array}{l} x', y' \text{ fresh} \quad H(x) = \iota \\ H(\iota) = _ \mathbf{obj} \ \{ _ \overline{M} \ \mathbf{method} \ m(y) \ \{ t \} \} \end{array}}{H; E[x.m(v)] \rightsquigarrow H, x' \mapsto \iota, y' \mapsto v; E[t[\mathbf{self} = x'][y = y']]}$	$\frac{\text{(R-SPAWN)} \quad \iota, i \text{ fresh} \quad x \notin \text{dom}(H) \quad H' = H, x \mapsto \iota, \iota \mapsto \mathbf{chan} \ \{ i, \emptyset \}}{H; E[\mathbf{spawn} \ (x) \ \{ t \}] \rightsquigarrow H'; E[\iota] \ t}$	$\frac{\text{(R-RECV)} \quad H(\iota) = \mathbf{chan} \ \{ i, \iota' \} \quad H' = H[\iota \mapsto \mathbf{chan} \ \{ i, \emptyset \}]}{H; E[\leftarrow \iota] \rightsquigarrow H'; E[\iota']}$		
$\frac{\text{(R-SENDERBLOCK)} \quad \begin{array}{l} H(\iota) = \mathbf{chan} \ \{ _ , \emptyset \} \quad \text{OkSend}(H, v) \\ i \text{ fresh} \quad H' = H[\iota \mapsto \mathbf{chan} \ \{ i, v \}] \end{array}}{H; E[\iota \leftarrow v] \rightsquigarrow H'; E[\blacksquare_i \ \iota]}$	$\frac{\text{(R-SENDERUNBLOCK)} \quad \begin{array}{l} H(\iota) = \mathbf{chan} \ \{ i', v \} \\ v = \emptyset \vee i \neq i' \end{array}}{H; E[\blacksquare_i \ \iota] \rightsquigarrow H; E[\iota]}$	$\frac{\text{(REFCHECK)} \quad \begin{array}{l} H(\iota) = K' \ \mathbf{obj} \ \{ _ _ \} \\ \text{OkField}(K, K') \end{array}}{\text{OkRef}(H, K, \iota)}$	$\frac{\text{(HELPER-OKSEND)} \quad \begin{array}{l} H(\iota) = K \ \mathbf{obj} \ \{ _ _ \} \\ \neg \exists \iota' \in \text{ROG}(H, \iota). \text{isLocal}(H, \iota') \end{array}}{\text{OkSend}(H, \iota)}$	$\frac{\text{(HELPER-OKFIELD)} \quad \begin{array}{l} K_3 = \text{Max}(K_1) \\ \forall K_4 \in K_2. K_3 \leq K_4 \end{array}}{\text{OkField}(K_1, K_2)}$

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Dalarna Dynamic Semantics

(R-UPDATE)

$$\begin{array}{c}
 H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ \overline{f = v} \ f = v' \ \overline{M} \} \\
 \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\
 H' = H[\iota \mapsto K \mathbf{obj} \{ \overline{f = v} \ f = v' \ \overline{M} \}] \\
 \hline
 H; E[x.f = v] \rightsquigarrow H'; E[v']
 \end{array}$$

(REFCHECK)

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 H(\iota) = K' \mathbf{obj} \{ _ _ \} \\
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 \end{array}$$

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$$\begin{array}{c}
 K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \\
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 K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}} \\
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Dalarna Dynamic Semantics

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Dalarna Dynamic Semantics

(R-UPDATE)

$$H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ \overline{f = v} \ f = v' \ \overline{M} \}$$

$$\neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v)$$

$$H' = H[\iota \mapsto K \mathbf{obj} \{ \overline{f = v} \ f = v' \ \overline{M} \}]$$

$$H; E[x.f = v] \rightsquigarrow H'; E[v']$$

(REFCHECK)

$$H(\iota) = K' \mathbf{obj} \{ _ _ \}$$

$$\text{OkField}(K, K')$$

$$\text{OkRef}(H, K, \iota)$$

(HELPER-OKFIELD)

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Dalarna Dynamic Semantics

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Dalarna Dynamic Semantics

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Dalarna Dynamic Semantics

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$$\text{OkField}(K_1, K_2)$$

Dalarna Dynamic Semantics

(R-UPDATE)

$$\begin{array}{c}
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 \end{array}$$

(REFCHECK)

$$\begin{array}{c}
 H(\iota) = K' \mathbf{obj} \{ _ _ \} \\
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$$K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1)$$

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$$\text{OkField}(K_1, K_2)$$

Dalarna Dynamic Semantics

(R-UPDATE)

$$\frac{\begin{array}{l} H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ \overline{f = v} \, f = v' \, \overline{M} \} \\ \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\ H' = H[\iota \mapsto K \mathbf{obj} \{ \overline{f = v} \, f = v' \, \overline{M} \}] \end{array}}{H; E[x.f = v] \rightsquigarrow H'; E[v']}$$

(REFCHECK)

$$\frac{\begin{array}{l} H(\iota) = K' \mathbf{obj} \{ _ _ \} \\ \text{OkField}(K, K') \end{array}}{\text{OkRef}(H, K, \iota)}$$

(HELPER-OKFIELD)

$$\begin{array}{l} K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \\ K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \\ K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}} \end{array}$$

$$\text{OkField}(K_1, K_2)$$

$$\text{local, iso} = \text{Max}(\text{local iso}), \text{Min}(\text{local iso})$$

$$\text{local, imm} = \text{Max}(\text{local imm}), \text{Min}(\text{local imm})$$

$$\text{local} \leq \text{local} \quad \text{iso} \leq \text{imm}$$

$$\text{OkField}(\text{local iso}, \text{local imm})$$

Dalarna Dynamic Semantics

(R-UPDATE)

$$\frac{\begin{array}{l} H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ \overline{f = v} \, f = v' \, \overline{M} \} \\ \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\ H' = H[\iota \mapsto K \mathbf{obj} \{ \overline{f = v} \, f = v' \, \overline{M} \}] \end{array}}{H; E[x.f = v] \rightsquigarrow H'; E[v']}$$

(REFCHECK)

$$\frac{\begin{array}{l} H(\iota) = K' \mathbf{obj} \{ _ _ \} \\ \text{OkField}(K, K') \end{array}}{\text{OkRef}(H, K, \iota)}$$

(HELPER-OKFIELD)

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$$\text{local, iso} = \text{Max}(\text{local iso}), \text{Min}(\text{local iso})$$

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Dalarna Dynamic Semantics

(R-UPDATE)

$$\begin{array}{c}
 H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ \overline{f = v} \ f = v' \ \overline{M} \} \\
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 H(\iota) = K' \mathbf{obj} \{ _ _ \} \\
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Dalarna Dynamic Semantics

(R-UPDATE)

$$\frac{\begin{array}{l} H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ \overline{f = v} \ f = v' \ \overline{M} \} \\ \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\ H' = H[\iota \mapsto K \mathbf{obj} \{ \overline{f = v} \ f = v' \ \overline{M} \}] \end{array}}{H; E[x.f = v] \rightsquigarrow H'; E[v']}$$

(REFCHECK)

$$\frac{\begin{array}{l} H(\iota) = K' \mathbf{obj} \{ _ _ \} \\ \text{OkField}(K, K') \end{array}}{\text{OkRef}(H, K, \iota)}$$

(HELPER-OKFIELD)

$$K_3^{Max}, K_3^{Min} = \text{Max}(K_1), \text{Min}(K_1)$$

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Dalarna Dynamic Semantics

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(REFCHECK)

$$\frac{\begin{array}{l} H(\iota) = K' \mathbf{obj} \{ _ _ \} \\ \text{OkField}(K, K') \end{array}}{\text{OkRef}(H, K, \iota)}$$

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$$\text{local} \leq \text{local} \quad \text{iso} \leq \text{imm}$$

$$\text{OkField}(\text{local iso}, \text{local imm})$$

$$\text{unsafe} \leq \text{local} \leq \text{iso} \leq \text{imm}$$

Dalarna Dynamic Semantics

(R-UPDATE)

$$\frac{\begin{array}{l} H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ \overline{f = v} \ f = v' \ \overline{M} \} \\ \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\ H' = H[\iota \mapsto K \mathbf{obj} \{ \overline{f = v} \ f = v' \ \overline{M} \}] \end{array}}{H; E[x.f = v] \rightsquigarrow H'; E[v']}$$

(REFCHECK)

$$\frac{\begin{array}{l} H(\iota) = K' \mathbf{obj} \{ _ _ \} \\ \text{OkField}(K, K') \end{array}}{\text{OkRef}(H, K, \iota)}$$

(HELPER-OKFIELD)

$$K_3^{Max}, K_3^{Min} = \text{Max}(K_1), \text{Min}(K_1)$$

$$K_4^{Max}, K_4^{Min} = \text{Max}(K_2), \text{Min}(K_2)$$

$$K_3^{Max} \leq K_4^{Max} \quad K_3^{Min} \leq K_4^{Min}$$

$$\text{OkField}(K_1, K_2)$$

$$\text{local, iso} = \text{Max}(\text{local iso}), \text{Min}(\text{local iso})$$

$$\text{local, imm} = \text{Max}(\text{local imm}), \text{Min}(\text{local imm})$$

$$\text{local} \leq \text{local} \quad \text{iso} \leq \text{imm}$$



$$\text{OkField}(\text{local iso}, \text{local imm})$$

$$\text{unsafe} \leq \text{local} \leq \text{iso} \leq \text{imm}$$

Dalarna Dynamic Semantics

$Configuration \quad Cfg ::= H, \bar{T}$
 $Thread \quad T ::= t \mid Err$
 $Thread\ Err \quad Err ::= Err_N \mid Err_A \mid Err_P \mid Err_C$

$\frac{(E-NoSuchField) \quad H(x.f) = \perp}{H; E[x.f] \rightsquigarrow H; Err_N}$	$\frac{(E-NoSuchMethod) \quad H(x) = \iota \quad H(\iota) = _ \mathbf{obj} \{ _ \bar{M} \} \quad m \notin names(\bar{M})}{H; E[x.m(v)] \rightsquigarrow H; Err_N}$	$\frac{(E-NoSuchFieldAssign) \quad H(x.f) = \perp}{H; E[x.f = v] \rightsquigarrow H; Err_N}$	$\frac{(E-SendBadTargetOrArgument) \quad H(\iota) = _ \mathbf{obj} \{ _ _ \} \vee H(\iota') = \mathbf{chan} \{ _ \}}{H; E[\iota \leftarrow \iota'], T \rightsquigarrow H; Err_N}$	
$\frac{(E-RecvBadTarget) \quad H(\iota) = _ \mathbf{obj} \{ _ _ \}}{H; E[\leftarrow \iota] \rightsquigarrow H; Err_N}$	$\frac{(E-CastError) \quad H(\iota) = K' \mathbf{obj} \{ _ _ \} \quad K' \neq K}{H; E[(K) \iota] \rightsquigarrow H; Err_C}$	$\frac{(E-AbsentVar) \quad H(x) = \top}{H; E[x] \rightsquigarrow H; Err_A}$	$\frac{(E-Consume) \quad H(x) = \top}{H; E[\mathbf{consume} \ x] \rightsquigarrow H; Err_A}$	$\frac{(E-AbsentTarget) \quad H(x) = \top}{H; E[x.m(v)] \rightsquigarrow H; Err_A}$
$\frac{(E-AbsentTargetAccess) \quad H(x) = \top}{H; E[x.f] \rightsquigarrow H; Err_A}$	$\frac{(E-AbsentFieldAssign) \quad H(x) = \top}{H; E[x.f = v] \rightsquigarrow H; Err_A}$	$\frac{(E-AbsentCopyTarget) \quad H(x) = \top}{H; E[K \mathbf{copy} \ x] \rightsquigarrow H; Err_A}$	$\frac{(E-BadFieldAssign) \quad H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ _ _ \} \quad isImm(H, \iota) \vee \neg OkRef(H, K, v)}{H; E[x.f = v] \rightsquigarrow H; Err_P}$	
$\frac{(E-BadInstantiation) \quad \neg \forall v \in \bar{v}. OkRef(H, K, v)}{H; E[K \mathbf{obj} \{ \overline{f = v \bar{M}} \}] \rightsquigarrow H; Err_P}$	$\frac{(E-SendingLocal) \quad H(v) = \mathbf{chan} \{ _ \} \quad \neg OkSend(H, \iota)}{H; E[v \leftarrow \iota], T \rightsquigarrow H; Err_P}$	$\frac{(E-AliasIso) \quad H(x) = \iota \quad isIso(H, \iota)}{H; E[x] \rightsquigarrow H; Err_P}$	$\frac{(E-IsoField) \quad H(x.f) = \iota \quad isIso(H, \iota)}{H; E[x.f] \rightsquigarrow H; Err_P}$	

Dalarna Dynamic Semantics

Configuration $Cfg ::= H, \bar{T}$
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$\frac{(E\text{-}NoSuchField) \quad H(x.f) = \perp}{H; E[x.f] \rightsquigarrow H; Err_N}$	$\frac{(E\text{-}NoSuchMethod) \quad H(x) = \iota \quad H(\iota) = _ \mathbf{obj} \{ _ \overline{M} \} \quad m \notin names(\overline{M})}{H; E[x.m(v)] \rightsquigarrow H; Err_N}$	$\frac{(E\text{-}NoSuchFieldAssign) \quad H(x.f) = \perp}{H; E[x.f = v] \rightsquigarrow H; Err_N}$	$\frac{(E\text{-}SendBadTargetOrArgument) \quad H(\iota) = _ \mathbf{obj} \{ _ _ \} \vee H(\iota') = \mathbf{chan} \{ _ \}}{H; E[\iota \leftarrow \iota'], T \rightsquigarrow H; Err_N}$	
$\frac{(E\text{-}RecvBadTarget) \quad H(\iota) = _ \mathbf{obj} \{ _ _ \}}{H; E[\leftarrow \iota] \rightsquigarrow H; Err_N}$	$\frac{(E\text{-}CastError) \quad H(\iota) = K' \mathbf{obj} \{ _ _ \} \quad K' \neq K}{H; E[(K) \iota] \rightsquigarrow H; Err_C}$	$\frac{(E\text{-}AbsentVar) \quad H(x) = \top}{H; E[x] \rightsquigarrow H; Err_A}$	$\frac{(E\text{-}Consume) \quad H(x) = \top}{H; E[\mathbf{consume} x] \rightsquigarrow H; Err_A}$	$\frac{(E\text{-}AbsentTarget) \quad H(x) = \top}{H; E[x.m(v)] \rightsquigarrow H; Err_A}$
$\frac{(E\text{-}AbsentTargetAccess) \quad H(x) = \top}{H; E[x.f] \rightsquigarrow H; Err_A}$	$\frac{(E\text{-}AbsentFieldAssign) \quad H(x) = \top}{H; E[x.f = v] \rightsquigarrow H; Err_A}$	$\frac{(E\text{-}AbsentCopyTarget) \quad H(x) = \top}{H; E[K \mathbf{copy} x] \rightsquigarrow H; Err_A}$	$\frac{(E\text{-}BadFieldAssign) \quad H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ _ _ \} \quad isImm(H, \iota) \vee \neg OkRef(H, K, v)}{H; E[x.f = v] \rightsquigarrow H; Err_P}$	
$\frac{(E\text{-}BadInstantiation) \quad \neg \forall v \in \overline{v}. OkRef(H, K, v)}{H; E[K \mathbf{obj} \{ \overline{f = v \overline{M}} \}] \rightsquigarrow H; Err_P}$	$\frac{(E\text{-}SendingLocal) \quad H(v) = \mathbf{chan} \{ _ \} \quad \neg OkSend(H, \iota)}{H; E[v \leftarrow \iota], T \rightsquigarrow H; Err_P}$	$\frac{(E\text{-}AliasIso) \quad H(x) = \iota \quad isIso(H, \iota)}{H; E[x] \rightsquigarrow H; Err_P}$	$\frac{(E\text{-}IsoField) \quad H(x.f) = \iota \quad isIso(H, \iota)}{H; E[x.f] \rightsquigarrow H; Err_P}$	

Dalarna Dynamic Semantics

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$\frac{(E\text{-}NoSuchField) \quad H(x.f) = \perp}{H; E[x.f] \rightsquigarrow H; Err_N}$	$\frac{(E\text{-}NoSuchMethod) \quad H(x) = \iota \quad H(\iota) = _ \mathbf{obj} \{ _ \overline{M} \} \quad m \notin names(\overline{M})}{H; E[x.m(v)] \rightsquigarrow H; Err_N}$	$\frac{(E\text{-}NoSuchFieldAssign) \quad H(x.f) = \perp}{H; E[x.f = v] \rightsquigarrow H; Err_N}$	$\frac{(E\text{-}SendBadTargetOrArgument) \quad H(\iota) = _ \mathbf{obj} \{ _ _ \} \vee H(\iota') = \mathbf{chan} \{ _ \}}{H; E[\iota \leftarrow \iota'], T \rightsquigarrow H; Err_N}$	
$\frac{(E\text{-}RecvBadTarget) \quad H(\iota) = _ \mathbf{obj} \{ _ _ \}}{H; E[\leftarrow \iota] \rightsquigarrow H; Err_N}$	$\frac{(E\text{-}CastError) \quad H(\iota) = K' \mathbf{obj} \{ _ _ \} \quad K' \neq K}{H; E[(K) \iota] \rightsquigarrow H; Err_C}$	$\frac{(E\text{-}AbsentVar) \quad H(x) = \top}{H; E[x] \rightsquigarrow H; Err_A}$	$\frac{(E\text{-}Consume) \quad H(x) = \top}{H; E[\mathbf{consume} x] \rightsquigarrow H; Err_A}$	$\frac{(E\text{-}AbsentTarget) \quad H(x) = \top}{H; E[x.m(v)] \rightsquigarrow H; Err_A}$
$\frac{(E\text{-}AbsentTargetAccess) \quad H(x) = \top}{H; E[x.f] \rightsquigarrow H; Err_A}$	$\frac{(E\text{-}AbsentFieldAssign) \quad H(x) = \top}{H; E[x.f = v] \rightsquigarrow H; Err_A}$	$\frac{(E\text{-}AbsentCopyTarget) \quad H(x) = \top}{H; E[K \mathbf{copy} x] \rightsquigarrow H; Err_A}$	$\frac{(E\text{-}BadFieldAssign) \quad H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ _ _ \} \quad isImm(H, \iota) \vee \neg OkRef(H, K, v)}{H; E[x.f = v] \rightsquigarrow H; Err_P}$	
$\frac{(E\text{-}BadInstantiation) \quad \neg \forall v \in \overline{v}. OkRef(H, K, v)}{H; E[K \mathbf{obj} \{ \overline{f = v \overline{M}} \}] \rightsquigarrow H; Err_P}$	$\frac{(E\text{-}SendingLocal) \quad H(v) = \mathbf{chan} \{ _ \} \quad \neg OkSend(H, \iota)}{H; E[v \leftarrow \iota], T \rightsquigarrow H; Err_P}$	$\frac{(E\text{-}AliasIso) \quad H(x) = \iota \quad isIso(H, \iota)}{H; E[x] \rightsquigarrow H; Err_P}$	$\frac{(E\text{-}IsoField) \quad H(x.f) = \iota \quad isIso(H, \iota)}{H; E[x.f] \rightsquigarrow H; Err_P}$	

Dalarna Dynamic Semantics

Configuration $Cfg ::= H, \bar{T}$
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(E-NoSuchField) $\frac{H(x.f) = \perp}{H; E[x.f] \rightsquigarrow H; \text{Err}_N}$	(E-NoSuchMethod) $\frac{H(x) = \iota \quad H(\iota) = _ \mathbf{obj} \{ _ \overline{M} \} \quad m \notin \text{names}(\overline{M})}{H; E[x.m(v)] \rightsquigarrow H; \text{Err}_N}$	(E-NoSuchFieldAssign) $\frac{H(x.f) = \perp}{H; E[x.f = v] \rightsquigarrow H; \text{Err}_N}$	(E-SendBadTargetOrArgument) $\frac{H(\iota) = _ \mathbf{obj} \{ _ _ \} \vee H(\iota') = \mathbf{chan} \{ _ \}}{H; E[\iota \leftarrow \iota'], T \rightsquigarrow H; \text{Err}_N}$	
(E-RecvBadTarget) $\frac{H(\iota) = _ \mathbf{obj} \{ _ _ \}}{H; E[\leftarrow \iota] \rightsquigarrow H; \text{Err}_N}$	(E-CastError) $\frac{H(\iota) = K' \mathbf{obj} \{ _ _ \} \quad K' \neq K}{H; E[(K) \iota] \rightsquigarrow H; \text{Err}_C}$	(E-AbsentVar) $\frac{H(x) = \top}{H; E[x] \rightsquigarrow H; \text{Err}_A}$	(E-Consume) $\frac{H(x) = \top}{H; E[\mathbf{consume} x] \rightsquigarrow H; \text{Err}_A}$	(E-AbsentTarget) $\frac{H(x) = \top}{H; E[x.m(v)] \rightsquigarrow H; \text{Err}_A}$
(E-AbsentTargetAccess) $\frac{H(x) = \top}{H; E[x.f] \rightsquigarrow H; \text{Err}_A}$	(E-AbsentFieldAssign) $\frac{H(x) = \top}{H; E[x.f = v] \rightsquigarrow H; \text{Err}_A}$	(E-AbsentCopyTarget) $\frac{H(x) = \top}{H; E[K \mathbf{copy} x] \rightsquigarrow H; \text{Err}_A}$	(E-BadFieldAssign) $\frac{H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ _ _ \} \quad \text{isImm}(H, \iota) \vee \neg \text{OkRef}(H, K, v)}{H; E[x.f = v] \rightsquigarrow H; \text{Err}_P}$	
(E-BadInstantiation) $\frac{\neg \forall v \in \overline{v}. \text{OkRef}(H, K, v)}{H; E[K \mathbf{obj} \{ \overline{f = v \overline{M}} \}] \rightsquigarrow H; \text{Err}_P}$	(E-SendingLocal) $\frac{H(v) = \mathbf{chan} \{ _ \} \quad \neg \text{OkSend}(H, \iota)}{H; E[v \leftarrow \iota], T \rightsquigarrow H; \text{Err}_P}$	(E-AliasIso) $\frac{H(x) = \iota \quad \text{isIso}(H, \iota)}{H; E[x] \rightsquigarrow H; \text{Err}_P}$	(E-IsoField) $\frac{H(x.f) = \iota \quad \text{isIso}(H, \iota)}{H; E[x.f] \rightsquigarrow H; \text{Err}_P}$	

Dalarna Dynamic Semantics

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(E-NoSuchField)

$$\frac{H(x.f) = \perp}{H; E[x.f] \rightsquigarrow H; Err_N}$$

(E-NoSuchMethod)

$$\frac{H(x) = \iota \quad H(\iota) = _ \mathbf{obj} \{ _ \bar{M} \} \quad m \notin \text{names}(\bar{M})}{H; E[x.m(v)] \rightsquigarrow H; Err_N}$$

(E-NoSuchFieldAssign)

$$\frac{H(x.f) = \perp}{H; E[x.f = v] \rightsquigarrow H; Err_N}$$

(E-SendBadTargetOrArgument)

$$\frac{H(\iota) = _ \mathbf{obj} \{ _ _ \} \vee H(\iota') = \mathbf{chan} \{ _ \}}{H; E[\iota \leftarrow \iota'], T \rightsquigarrow H; Err_N}$$

(E-RecvBadTarget)

$$\frac{H(\iota) = _ \mathbf{obj} \{ _ _ \}}{H; E[\leftarrow \iota] \rightsquigarrow H; Err_N}$$

(E-CastError)

$$\frac{H(\iota) = K' \mathbf{obj} \{ _ _ \} \quad K' \neq K}{H; E[(K) \iota] \rightsquigarrow H; Err_C}$$

(E-AbsentVar)

$$\frac{H(x) = \top}{H; E[x] \rightsquigarrow H; Err_A}$$

(E-Consume)

$$\frac{H(x) = \top}{H; E[\mathbf{consume} x] \rightsquigarrow H; Err_A}$$

(E-AbsentTarget)

$$\frac{H(x) = \top}{H; E[x.m(v)] \rightsquigarrow H; Err_A}$$

(E-AbsentTargetAccess)

$$\frac{H(x) = \top}{H; E[x.f] \rightsquigarrow H; Err_A}$$

(E-AbsentFieldAssign)

$$\frac{H(x) = \top}{H; E[x.f = v] \rightsquigarrow H; Err_A}$$

(E-AbsentCopyTarget)

$$\frac{H(x) = \top}{H; E[K \mathbf{copy} x] \rightsquigarrow H; Err_A}$$

(E-BadFieldAssign)

$$\frac{H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ _ _ \} \quad \text{isImm}(H, \iota) \vee \neg \text{OkRef}(H, K, v)}{H; E[x.f = v] \rightsquigarrow H; Err_P}$$

(E-BadInstantiation)

$$\frac{\neg \forall v \in \bar{v}. \text{OkRef}(H, K, v)}{H; E[K \mathbf{obj} \{ \overline{f = v \bar{M}} \}] \rightsquigarrow H; Err_P}$$

(E-SendingLocal)

$$\frac{H(v) = \mathbf{chan} \{ _ \} \quad \neg \text{OkSend}(H, \iota)}{H; E[v \leftarrow \iota], T \rightsquigarrow H; Err_P}$$

(E-AliasIso)

$$\frac{H(x) = \iota \quad \text{isIso}(H, \iota)}{H; E[x] \rightsquigarrow H; Err_P}$$

(E-IsoField)

$$\frac{H(x.f) = \iota \quad \text{isIso}(H, \iota)}{H; E[x.f] \rightsquigarrow H; Err_P}$$

Dalarna Properties

Theorem 4.4 (Dalarna is Data-Race Free Modulo Unsafe Objects). *Any data race in Dalarna directly or indirectly involves an unsafe object. A data race is defined as a read/write or write/write access to an object from different threads without any interleaving synchronisation, which in our case means a transfer of the object from one thread to the other.*

Theorem 4.5 (Preservation). *Given a well-formed configuration $\Gamma; H \vdash t \bar{T}$ and $H; t \bar{T} \rightsquigarrow H'; \bar{T}' \bar{T}$ then, there exists a Γ' s.t. $\Gamma' \supseteq \Gamma$ and $\Gamma'; H' \vdash \bar{T}' \bar{T}$*

Theorem 4.6 (Progress). *Given a well-formed configuration $\Gamma; H \vdash \bar{T}$, then either $\Gamma; H \vdash \bar{T}$ is a terminal configuration (see below) or $H; \bar{T} \rightsquigarrow H'; \bar{T}'$.*

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Future Work

- Addition of capability-based gradual type system to statically reject ill-capable programs

```
ls = iso object List {...}  
ls.append(unsafe object 0 {...} )   Statically rejected
```

- Improve the prototype and benchmark performance

Transient Typechecks Are (Almost) Free. [ECOOP 2019](#)

Richard Roberts, Stefan Marr, Michael Homer, James Noble

Conclusions



- Dalarna programming model
 - guarantees data race free programs
 - safe shared memory without deep copying
 - can be embedded in existing languages (Grace)

Image Source



"Dala Horse" Art Prints by JoniandCo | Redbubble

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